

Introduction

- Goal: Design an iterative IA algorithm that achieves high DoF for the case of a constant, K user, fully symmetric, generic, MIMO interference channel, assuming perfect CSIT.
- Contributions:
 - Design of vector space IA scheme is recast as a nonconvex quadratic feasibility problem.
 - ii. An extension of the **Feasible Point Pursuit Successive Convex Approximation** (FPP-SCA) [Mehanna *et al.*, 2015] algorithm is proposed to obtain high-quality suboptimal solutions.

System Model

- K users, each with N Tx-Rx antenna pairs, transmitting d streams.
- Estimated signal at k^{th} receiver given by

 $\hat{\mathbf{s}}_k = \mathbf{U}_k^T \mathbf{H}_{kk} \mathbf{V}_k \mathbf{s}_k + \sum_{j=1, j \neq k}^K \mathbf{U}_k^T \mathbf{H}_{kj} \mathbf{V}_j \mathbf{s}_j + \mathbf{U}_k^T \mathbf{n}_k$

- $\mathbf{s}_k \in \mathbb{R}^{d \times 1}$ is the transmitted symbol vector of user k.
- $\mathbf{V}_k \in \mathbb{R}^{N \times d}$ is the linear precoding matrix at k^{th} transmitter.
- $\mathbf{H}_{ki} \in \mathbb{R}^{N \times N}$ is the channel coefficient matrix between transmitter and receiver.
- $\mathbf{n}_k \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_N)$
- $\mathbf{U}_k \in \mathbb{R}^{N \times d}$ is the post processing matrix at k^{th} receiver.

Problem Description

For generic channels, problem is equivalent to

find $\{\mathbf{U}_k, \mathbf{V}_k\}_{k=1}^K$

s.t.

 $\{\mathbf{U}_k,\mathbf{V}_k\}_{k=1}^K$ (1a) $\mathbf{U}_{k}^{T}\mathbf{H}_{kj}\mathbf{V}_{j} = \mathbf{0}, \forall \ j \neq k \in \mathcal{K}$ (1b) $\operatorname{Rank}(\mathbf{U}_k) = \operatorname{Rank}(\mathbf{V}_k) = d, \forall k \in \mathcal{K}$ (1c)

Nonconvex, no polynomial-time algorithm known

Prior Art

- **Leakage Minimization:** distributed algorithm, minimize leakage interference via alternating minimization, low complexity. [Gomadam et *al.*, 2008].
- **MAX SINR:** distributed algorithm, maximize SINR at receivers, low complexity. [Gomadam et al., 2008].
- Rank Constrained Rank Minimization: centralized algorithm, minimize nuclear norm heuristic via alternating SDPs, high complexity. [Papailiopoulos *et al.*, 2011]



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Problem Reformulation

example $\mathbf{u}_{k}^{\{l\}T}\mathbf{H}_{k,j}\mathbf{v}_{j}^{\{m\}}=0,$

 $\Leftrightarrow \begin{bmatrix} \mathbf{u}_k^{\{l\}T} & \mathbf{v}_j^{\{m\}T} \end{bmatrix} \begin{bmatrix} \mathbf{0}_N \\ \mathbf{H} \end{bmatrix}$

 $\Leftrightarrow (\mathbf{x}_{k,j}^{\{l,m\}})^T \mathbf{G}_{k,j} (\mathbf{x}_{k,j}^{\{l,m\}})$

Overall, we have the problem

 $\inf_{\{\mathbf{U}_k,\mathbf{V}_k\}_{k=1}^K} \{\mathbf{U}_k,\mathbf{V}_k\}_{k=1}^K$ s.t. $\mathbf{x}_{k,j}^{\{l,m\}} = \left[\mathbf{u}_k^{\{l\}}\right]$ (Zero forcing constraints) $(\mathbf{x}_{k,i}^{\{l,m\}})^T \mathbf{G}_{k,j}$ $\|\mathbf{u}_k^{\{l\}}\|_2^2 \ge 1, \mathbf{)}$ (Unit norm constraints) $\|\mathbf{v}_{k}^{\{l\}}\|_{2}^{2} \ge 1$ $\mathbf{y}_k^{\{l,m\}} = \left[\mathbf{u}_k^{\{l\}}\right]$

(Orthogonality constraints)

 $\mathbf{z}_{k}^{\{l,m\}} = \left[\mathbf{v}_{k}^{\{l\}}\right]$ $(\mathbf{z}_{k}^{\{l,m\}})^{T}\mathbf{E}(\mathbf{z}_{k}^{\{l,m\}})$

FPP - SCA Algorithm

$$\begin{split} P_{i} & \min_{\{\mathbf{U}_{k},\mathbf{V}_{k}\}_{k=1}^{K}} \sum_{\substack{k,j \in \mathcal{K}, j \neq k, \\ m,l \in \{1, \cdots, d\}}} (s_{k,j}^{\{l,m\}}(1) + s_{k,j}^{\{l,m\}}(2)) \\ & + \sum_{\substack{k \in \mathcal{K}, \\ l \in \{1, \cdots, d\}}} (s_{k}^{\{l\}}(1) + s_{k}^{\{l\}}(2)) \\ & + \sum_{\substack{k \in \mathcal{K}, \\ m \neq l \in \{1, \cdots, d\}}} (s_{k}^{\{l,m\}}(1) + s_{k}^{\{l,m\}}(2) + s_{k}^{\{l,m\}}(3) + s_{k}^{\{l,m\}}(4)) \\ & \times_{k,j}^{\{l,m\}} = \left[\mathbf{u}_{k}^{\{l\}T} \cdot \mathbf{v}_{j}^{\{m\}T}\right]^{T}, \\ \text{s.t.} & \times_{k,j}^{\{l,m\}} \in \mathcal{C}_{2F}(\mathbf{u}_{k}^{\{l\}}, \mathbf{v}_{j}^{\{m\}}), \\ & s_{k,j}^{\{l,m\}}(1) \geq 0, \ s_{k,j}^{\{l,m\}}(2) \geq 0 \\ & \mathbf{u}_{k}^{\{l\}} \in \mathcal{C}_{\mathrm{UN}}(\mathbf{u}_{k}^{\{l\}}), \\ & \mathbf{v}_{k}^{\{l\}} \in \mathcal{C}_{\mathrm{UN}}(\mathbf{v}_{k}^{\{l\}}), \\ & \mathbf{v}_{k}^{\{l\}} \in \mathcal{C}_{\mathrm{UN}}(\mathbf{v}_{k}^{\{l\}}), \\ & s_{k}^{\{l,m\}}(1) \geq 0, \ s_{k}^{\{l,m\}}(2) \geq 0 \\ & \mathbf{v}_{k}^{\{l,m\}} = \left[\mathbf{u}_{k}^{\{l\}T} \cdot \mathbf{u}_{k}^{\{m\}T}\right]^{T}, \\ & \mathbf{v}_{k}^{\{l,m\}} \in \mathcal{C}_{\mathrm{ORTH}}(\mathbf{u}_{k}^{\{l\}}, \mathbf{u}_{k}^{\{m\}}), \\ & s_{k}^{\{l,m\}} \in \left[\mathbf{v}_{k}^{\{l,m\}} \cdot \mathbf{v}_{k}^{\{m\}T}\right]^{T}, \\ & \mathbf{v}_{k}^{\{l,m\}} \in \mathcal{C}_{\mathrm{ORTH}}(\mathbf{v}_{k}^{\{l\}}, \mathbf{v}_{k}^{\{m\}}), \\ & \mathbf{v}_{k}^{\{l,m\}} \in \mathcal{C}_{\mathrm{ORTH}}(\mathbf{v}_{k}^{\{l\}}, \mathbf{v}_{k}^{\{m\}}), \\ & \mathbf{v}_{k}^{\{l,m\}} \in \mathcal{C}_{\mathrm{ORTH}}(\mathbf{v}_{k}^{\{l\}}, \mathbf{v}_{k}^{\{m\}}), \\ & \forall l \neq m \in \{1, \cdots, d\}, \\ & k \in \mathcal{K} \\ & \mathsf{interms} = \left[\mathbf{v}_{k}^{\{l,m\}} \cdot \mathbf{v}_{k}^{\{m\}T}\right]^{T}, \\ & \mathbf{v}_{k}^{\{l,m\}} \in \mathcal{C}_{\mathrm{ORTH}}(\mathbf{v}_{k}^{\{l\}}, \mathbf{v}_{k}^{\{m\}}), \\ & \forall l \neq m \in \{1, \cdots, d\}, \\ & k \in \mathcal{K} \\ & \mathsf{interms} = \left[\mathbf{v}_{k}^{\{l,m\}} \cdot \mathbf{v}_{k}^{\{m\}}\right]^{T}, \\ & \forall l \neq m \in \{1, \cdots, d\}, \\ & k \in \mathcal{K} \\ & \mathsf{interms} = \left[\mathbf{v}_{k}^{\{l,m\}} \cdot \mathbf{v}_{k}^{\{m\}}\right]^{T}, \\ & \mathsf{interms} = \left[\mathbf{v}_{k}^{\{m\}} \cdot \mathbf{v}_{k}^{\{m\}}\right]^{T}, \\ & \mathsf{inte$$

Reformulate each constraint as a quadratic equality constraint. For

$$\begin{array}{l} \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \\ \begin{array}{l} \mathbf{W}_{k \times N} & \mathbf{H}_{k, j} \\ \mathbf{H}_{k, j}^{T} & \mathbf{0}_{N \times N} \end{array} \right] \begin{bmatrix} \mathbf{u}_{k}^{\{l\}} \\ \mathbf{v}_{j}^{\{m\}} \end{bmatrix} = 0, \\ \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \\ \begin{array}{l} \uparrow^{n} \end{array} \right) = 0, \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \\ \begin{array}{l} \mathbf{n} \end{array} \right) = 0, \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \\ \begin{array}{l} \mathbf{n} \end{array} \right) \\ \begin{array}{l} \uparrow^{T} & \mathbf{v}_{j}^{\{m\}T} \end{bmatrix}^{T}, \\ \begin{array}{l} \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \end{array} \right) \\ \begin{array}{l} \forall \ l, m \in \{1, 2, \cdots, d\}, \ j \neq k \in \mathcal{K} \end{array} \right) \\ \begin{array}{l} \downarrow \ j \neq k \in \mathcal{K} \end{array}$$

$$\begin{aligned} \|\mathbf{u}_{k}^{k}\|_{2} &\geq 1 \\ \|\mathbf{v}_{k}^{\{l\}}\|_{2}^{2} &\geq 1 \end{aligned} \right\} \forall l \in \{1, \cdots, d\}, \ k \in \mathcal{K} \\ \mathbf{y}_{k}^{\{l,m\}} &= \begin{bmatrix} \mathbf{u}_{k}^{\{l\}T} & \mathbf{u}_{k}^{\{m\}T} \end{bmatrix}^{T}, \\ (\mathbf{y}_{k}^{\{l,m\}})^{T} \mathbf{E}(\mathbf{y}_{k}^{\{l,m\}}) &\geq 0 \end{aligned} \right\} \forall l \neq m \in \{1, \cdots, d\}, \\ \mathbf{z}_{k}^{\{l,m\}} &= \begin{bmatrix} \mathbf{v}_{k}^{\{l\}T} & \mathbf{v}_{k}^{\{m\}T} \end{bmatrix}^{T}, \\ (\mathbf{z}_{k}^{\{l,m\}})^{T} \mathbf{E}(\mathbf{z}_{k}^{\{l,m\}}) &\geq 0 \end{aligned} \right\} \forall l \neq m \in \{1, \cdots, d\}, \\ (\mathbf{z}_{k}^{\{l,m\}})^{T} \mathbf{E}(\mathbf{z}_{k}^{\{l,m\}}) &\geq 0 \end{aligned}$$

Main Idea: Use convex inner approximation of feasible set about random starting point and apply positive slacks to maintain feasibility. Impose ℓ_1 penalty on slacks to minimize constraint violations.

Overall Algorithm $\begin{cases} \mathbf{Step 0:} \\ \{\mathbf{p}_{k,j}^{\{l,m\}(0)}\}, \end{cases}$ Randomly generate initialization points $\{\mathbf{q}_{k}^{\{l\}(0)}, \mathbf{r}_{k}^{\{l\}(0)}\}, \{\mathbf{t}_{k}^{\{l,m\}(0)}, \mathbf{w}_{k}^{\{l,m\}(0)}\}\$ for each pair of constraints in (P_0) . **Step i:** Solve the problem (P_i) to obtain a set of solutions $\{\mathbf{U}_{k}^{(i)}, \mathbf{V}_{k}^{(i)}\}_{k=1}^{K}$. Update $\{\mathbf{p}_{k,j}^{\{l,m\}(i+1)} := \begin{bmatrix} \mathbf{u}_k^{\{l\}(i)T} & \mathbf{v}_j^{\{m\}(i)T} \end{bmatrix}^T \};$ $\{\mathbf{q}_{k}^{\{l\}(i+1)}\} := \mathbf{u}_{k}^{\{l\}(i)}; \{\mathbf{r}_{k}^{\{l\}(i+1)}, \} := \mathbf{v}_{k}^{\{l\}(i)};$ $\{\mathbf{t}_{k}^{\{l,m\}(i+1)} := \begin{bmatrix} \mathbf{u}_{k}^{\{l\}(i)^{T}} & \mathbf{u}_{k}^{\{m\}(i)^{T}} \end{bmatrix}^{T} \};$ $\{\mathbf{w}_{k}^{\{l,m\}(i+1)} := \begin{bmatrix} \mathbf{v}_{k}^{\{l\}(i)^{T}} & \mathbf{v}_{k}^{\{m\}(i)^{T}} \end{bmatrix}^{T} \};$ Repeat until feasibilty achieved or maximum number of iterations Generates a monotonically non-increasing, convergent cost sequence. If cost function converges to zero, a solution of problem (1) will be obtained. **Simulation Results** FPP-SCA run for a maximum of 120 iterations, with a maximum of 5 restarts. Leakage minimization and Max SINR run for maximum of 10⁴ iterations. Fig 1: DoF for a $(3 \times 3, 1)^5$ MIMO IC Alt. Min. -- - OFDM DoF UB Performance of FPP-SCA very close to theoretical upper bound. Other algorithms significantly worse in comparison. Fig 2: DoF for a $(6 \times 6, 2)^5$ MIMO IC DoF vs P 🕒 😑 🗕 Max SINF ---- OFDM FPP-SCA achieves significantly higher DoF in comparison to other algorithms. **Future Work:** Low complexity version of FPP-SCA.

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Extension to more general channel settings, e.g., SISO IC with generic channel extensions.